9. TCP/IP

A. What is TCP/IP?

- TCP/IP is a set of protocols developed to allow cooperating computers to share resources across a network.
- **TCP** stands for “Transmission Control Protocol”
- **IP** stands for “Internet Protocol”
- They are **Transport layer and Network layer** protocols respectively of the protocol suite.
- The most well known network that adopted TCP/IP is **Internet** – the biggest WAN in the world.
What is a protocol?

• A protocol is a collection of rules and procedures for two computers to exchange information
• Protocol also defines the format of data that is being exchanged
Why TCP/IP is so popular?

- TCP/IP was developed very early
- Technologies were widely discussed and circulated in documents called “Request for Comments” (RFC) – free of charge
- Supported by UNIX operating system
TCP/IP Model

- Because TCP/IP was developed earlier than the OSI 7-layer mode, it does not have 7 layers but only 4 layers

**TCP/IP Protocol Suite**

- FTP, SMTP, Telnet, HTTP,…
- TCP, UDP
- IP, ARP, ICMP
- Network Interface

**OSI 7-layer**

1. Physical Layer
2. Data Link Layer
3. Network Layer
4. Transport Layer
5. Session Layer
6. Presentation Layer
7. Application Layer
• Application layer protocols define the rules when implementing specific network applications
• Rely on the underlying layers to provide accurate and efficient data delivery
• Typical protocols:
  • FTP – File Transfer Protocol
    • For file transfer
  • Telnet – Remote terminal protocol
    • For remote login on any other computer on the network
  • SMTP – Simple Mail Transfer Protocol
    • For mail transfer
  • HTTP – Hypertext Transfer Protocol
    • For Web browsing
TCP/IP is built on “connectionless” technology, each datagram finds its own way to its destination.

Transport Layer protocols define the rules of:
- Dividing a chunk of data into segments
- Reassemble segments into the original chunk

Typical protocols:
- **TCP** – Transmission Control Protocol
  - Provide further the functions such as reordering and data resend
- **UDP** – User Datagram Service
  - Use when the message to be sent fit exactly into a datagram
  - Use also when a more simplified data format is required
Network layer protocols define the rules of how to find the routes for a packet to the destination.

It only gives best effort delivery. Packets can be delayed, corrupted, lost, duplicated, out-of-order.

Typical protocols:

- **IP** – Internet Protocol
  - Provide packet delivery
- **ARP** – Address Resolution Protocol
  - Define the procedures of network address / MAC address translation
- **ICMP** – Internet Control Message Protocol
  - Define the procedures of error message transfer
Application Layer

- Application
- Transport
- Network
- Network Interface
B. Example: SMTP

Client

*SMTP*

- TCP
- IP, ARP, ICMP
- Network Interface

Virtual

SMTP Server

*SMTP*

- TCP
- IP, ARP, ICMP
- Network Interface

Actual
• The underlying layers have guaranteed accurate data delivery
• We need to make a lot agreements with the server in application layer before sending mail

1. Agree on how data is represented
   • Binary or ASCII
2. Ensure the right recipient
   • There may be 1000 users served by the server
3. Ensure the client has the right to send mail
   • Some clients are not welcome
4. How to tell the server it is the end of the message
   • All mail looks the same
• Example: SMTP

The following mail is to be sent:

Date: Fri, 18 Jan 02 13:26:31 EDT
From: enpklun@polyu.edu.hk
To: tchsun@eee.hku.hk
Subject: meeting

Let’s get together Monday at 1pm.
SMTP Server

access port 25 of server

Client

220 eee.hku.hk SMTP Service at 20 Jan 02 05:17:18 EDT

HELO polyu.edu.hk

250 eee.hku.hk – Hello, polyu.edu.hk

MAIL From:
<enpklun@polyu.edu.hk>

250 MAIL accepted

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9. TCP/IP
Client

SMTP Server

RCPT To: <tchsun@eee.hku.hk>

DATA

250 Recipient accepted

354 Start mail input; end with .

Date: Fri, 18 Jan 02 13:26:31 EDT
From: enpklun@polyu.edu.hk
To: tchsun@eee.hku.hk
Subject: meeting

Let’s get together Monday at 1pm.
• The agreement made in the SMTP protocol
  • All messages use normal text
    • All ASCII characters
  • The responses all begin with numbers
    • To indicate the status when receiving the command
  • Some words are reserved words
    • HELO, MAIL, RCPT…
  • Mail ends with a line that contains only a period

• The information passed with the SMTP messages
  • The recipient name
  • The sender name
  • The mail
C. Domain Name

• Every computer has a network address
  • e.g. 158.132.161.99
• To access a computer, we need to specify its network address
• Human beings are weak in memorizing numbers
• We prefer computer name or domain name
  • e.g. hkpu10.polyu.edu.hk
• Need a machine on the Internet to convert name to number
Domain name hierarchy

Example:

hkpu10.polyu.edu.hk

Computer name

- The domain within edu.hk
- One of the educational institutions in H.K.

Root domain name

.other examples:
- com – commercial company
- org – general organization
- net – major network centre
- gov – government org.
- mil – military group
- edu – education org.
• An organization needs to register its domain name
  • e.g. PolyU has registered its name to the domain of edu.hk
• Once a domain name is assigned, the organization is free to assign other names belong to its domain
  • e.g. we can have
    hkpu10.polyu.edu.hk
    smtp.polyu.edu.hk
    mail.polyu.edu.hk
Domain Name Server (DNS) of polyu.edu.hk

Where is www.yahoo.com?

usually UDP

Address of www.yahoo.com

Where is www.yahoo.com?

Address of yahoo.com?

DNS of Yahoo.com

Where is yahoo.com?

DNS of com

Address of the DNS of Yahoo.com

Client

DNS of Yahoo.com

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9. TCP/IP
• Nevertheless, such a complicated procedure needs not perform in most cases
• Client computers usually remember the answers that it got before
• It reduces the loading to the root DNS
• To further reduce loading, there can be many root DNS on the Internet
  • e.g. there are a few “com” root DNS
Transport Layer

Application

Transport

Network

Network Interface

Message

Segments

h M

h M

h M
D. TCP and UDP

TCP – Transmission Control Protocol

• TCP is a connection-oriented protocol
  • Does not mean it has a physical connection between sender and receiver
  • TCP provides the function to allow a connection virtually exists – also called virtual circuit

• TCP provides the functions:
  • Dividing a chunk of data into segments
  • Reassembly segments into the original chunk
  • Provide further the functions such as reordering and data resend

• Offering a reliable byte-stream delivery service
Dividing and Reassembly

TCP

Message

Source Port | Destination Port
---|---
Sequence Number
Acknowledgement Number
Checksum
Message Data
Sender

1

2

3

A1

A2

A3

Timeout

retransmit

Recipient

1

2

3
• A Typical Procedure
  • Sender
    • TCP divides a message into segments
    • Add sequence no.
    • Send the segments in sequence and wait for acknowledgement
    • If an acknowledgement for a segment is not received for a certain period of time, resend it until an acknowledgement is received
  • Recipient
    • When receiving segments, send the acknowledgement with correct number
    • Reassembly the segments back to the message
Port Multiplexing

- A computer may perform a number of network applications at the same time
  - FTP + SMTP + HTTP, etc.
- Each computer has only one network address, how can it serve so many applications at the same time?

⇒ by port multiplexing

Network add: 158.132.161.99
Well-known Port Numbers

• Some port numbers are reserved for some purposes
  • Port 21: FTP – file transfer
  • Port 25: SMTP – mail transfer
  • Port 23: TELNET – remote login
  • Port 80: HTTP – Web access

• These port numbers are well known to all computers in the network

• E.g. whenever a client access port 25 of the server, it means the client needs SMTP service
### SMTP Server

- Located by: network address + TCP port no.

<table>
<thead>
<tr>
<th>Source Port</th>
<th>Destination Port</th>
</tr>
</thead>
<tbody>
<tr>
<td>1357</td>
<td>25</td>
</tr>
</tbody>
</table>

- Sequence Number
- Acknowledgement Number
- Checksum
- Message Data
Client A

SMTP port = 1357

Client B

FTP port = 1361

SMTP + FTP Server

Network address: 158.132.161.99

SMTP port = 25

FTP port = 21
Network Layer

Application

Transport

Network

Network Interface

Message

Segments

Datagrams / Packets

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9. TCP/IP
E. Network Addresses and Subnets

- A header is added to each segment in the Network layer

<table>
<thead>
<tr>
<th>Total Length</th>
</tr>
</thead>
<tbody>
<tr>
<td>Time to Live</td>
</tr>
<tr>
<td>Source Address</td>
</tr>
<tr>
<td>Destination Address</td>
</tr>
</tbody>
</table>

Segment
• **Total Length** – Total length of a packet (up to 65535 bytes)
• **Time to Live** – How many times this packet can be routed on the network (up to 255)
• **Protocol** – The transport layer protocol that the packet belongs to
  • TCP: 6
  • UDP: 17
  • ICMP: 1
• **Source address** – the network address of the computer that sends the data
• **Destination address** – the network address of the computer that the data is sending to
• Each computer (host) must have a unique network address (or IP address for TCP/IP suite)
• Each IP address is 32-bit long (four bytes)
• The four-byte address is written out as a.b.c.d
  • e.g.

<table>
<thead>
<tr>
<th>Byte 1</th>
<th>Byte 2</th>
<th>Byte 3</th>
<th>Byte 4</th>
</tr>
</thead>
<tbody>
<tr>
<td>158</td>
<td>132</td>
<td>161</td>
<td>99</td>
</tr>
</tbody>
</table>

• IP addresses are hierarchical
  • network I.D. and host I.D.
• Each Network I.D. on the Internet needs to be registered to the Internet Assigned Number Authority
Class A – for very large network

1 bit  7 bits  24 bits

<table>
<thead>
<tr>
<th></th>
<th>Net I.D.</th>
<th>Host I.D.</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

• Only $2^7$ (63) networks can belong to this class
• Each network, there are $2^{24}$ hosts or computers
• Very few class A networks in the world
  • e.g. Arpanet – the earliest packet switched WAN (started 40 years ago)
Class B – for medium size network

- 2 bits
- 14 bits
- 16 bits

<table>
<thead>
<tr>
<th>1</th>
<th>0</th>
<th>Net I.D.</th>
<th>Host I.D.</th>
</tr>
</thead>
</table>

- $2^{14}$ (16384) networks can belong to this class
- Each network, there are $2^{16}$ (65536) hosts or computers
- Polyu’s address belongs to this group
  - e.g. 158.132.14.1

```
1001 1110
1000 0100
0000 1110
0000 0001
```

Network I.D. Host I.D.
Class C – for small network

- $2^{21}$ networks can belong to this class
- Each network, there are only $2^8$ (256) hosts or computers
Class D – for multicast network

- Packets are addressed to a multicast group
- Not often supported on Internet
Special Addresses

- **Host I.D. = all ‘1’s** ⇒ Directed broadcast
  “Broadcast to all hosts in the network or subnetwork”, not assigned
- **Host I.D. = all ‘0’s** ⇒ “This network”, not assigned
- **Network I.D. = 127** is reserved for loopback and diagnostic purposes, not assigned
- **Network I.D. + Host I.D. = all ‘1’s** ⇒ Limited broadcast
  “Broadcast to all hosts in the current network”, not assigned
Subnets

- A class B address can have 65536 hosts
- Difficult to manage
- Usually subdivide into a few small subnets
- Subnetting can also help to reduce broadcasting traffic

```
All traffic to 158.132.0.0
```

```
Each subnet 256 hosts
```

```
158.132.0.0
158.132.1.0
158.132.2.0
158.132.3.0
```

```
Total 65536 hosts
```
Subnet Mask

• How does the router know which subnet a packet should go?
• For each interface of the router, a subnet mask is provided to redefine which part of the address is Net ID and which part is Host ID
• Become classless addressing

A subnet mask: 255.255.255.0

1111 1111.1111 1111. 1111 1111. 0000 0000

‘1’s  Net ID                  ‘0’s  Host ID
A packet with destination address 158.132.1.10

Routing Table

<table>
<thead>
<tr>
<th>Subnet</th>
<th>S0</th>
<th>S1</th>
<th>S2</th>
</tr>
</thead>
<tbody>
<tr>
<td>158.132.1.0</td>
<td>158.132.2.0</td>
<td>158.132.3.0</td>
<td></td>
</tr>
<tr>
<td>Mask</td>
<td>255.255.255.0</td>
<td>255.255.255.0</td>
<td>255.255.255.0</td>
</tr>
</tbody>
</table>

158.132.1.10 1001 1110.1000 0100.0000 0001.0000 1010
AND 255.255.255.0 AND 1111 1111.1111 1111.1111 1111.0000 0000
158.132.1.0 1001 1110.1000 0100.0000 0001.0000 0000

Advantage: easy to compute
F. Routing

• How a packet finds its way to a computer in a network?
  • By using Routers
  • Routing is the selection of a path to guide a packet from the source to the destination
• Criteria in selecting a path may be:
  • Shortest path
  • Quickest path
  • Cheapest path
The red path is the shortest path.
• Each router has a table that records the estimated distance to all other routers
• If a router knows the entire network topology, the shortest path can be calculated
• To achieve this, routers broadcast Link State Advertisement to all other routers periodically
  • By means of routing protocol
• Each router knows the exact topology, and then calculates the shortest path
• In practice, it is not possible for a router to all paths. Only the nearer ones are kept
  • Hence can give wrong estimation
Host A
158.132.148.66
Default gateway: Router C

Router C

Router A

Routing Table
S0 160.64.124.0  Direct
255.255.255.0
S1 160.64.123.0  Direct
255.255.255.0

Subnet
158.132.166.0
Routing Table
S1 158.132.166.0  Direct
255.255.255.0
T1 160.64.0.0  Forward
255.255.0.0

Subnet
160.64.124.0

Subnet
160.64.123.0

Host B
160.64.123.98

Routing Table
S0 160.64.124.0  Direct
255.255.255.0
S1 160.64.123.0  Direct
255.255.255.0

Router B

Host A
158.132.148.66
Default gateway: Router C

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1. Host A wants to send a packet to Host B with address 160.64.123.98
2. Host A checks that 160.64.123.98 is not in the same network
3. Send packet to default gateway (Router C)
4. Default gateway finds that it cannot provide the best route for the packet, inform Host A to send the packet to Router A next time
5. Router C sends the packet to Router A
6. Router A checks from the table the packet should forward to Router B
7. Router B receives the packet and checks in its table the packet should directly deliver to subnet 160.64.123.0
8. Host B (160.64.123.98) receives the packet
Data Link and Physical Layers

Application

Transport

Network

Network Interface

Message

Segments

Packets

Frames

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9. TCP/IP
G. Ethernet Encapsulation and ARP

• An IP packet should be **encapsulated** into a frame for transmission by data link layer
• e.g. if **Ethernet (or IEEE 802.3)** is used:

<table>
<thead>
<tr>
<th>Preamble</th>
<th>Des. Add</th>
<th>Sour. Add</th>
<th>Length</th>
<th>IP Packet</th>
<th>FCS</th>
</tr>
</thead>
<tbody>
<tr>
<td>7 Bytes</td>
<td>1 Byte</td>
<td>2/6 Bytes</td>
<td>2/6 Bytes</td>
<td>2 Bytes</td>
<td>46 - 1500 Bytes</td>
</tr>
</tbody>
</table>

**IEEE 802.3 Frame**
• Only the hardware address (MAC address) is unique to a host
• Need to **convert** a network address to MAC address

Source IP = 158.132.148.66  
Destination IP = 158.132.148.132

Ethernet address = ?
**ARP – Address Resolution Protocol**

**Case 1**

1. Broadcast: Who has got IP address 158.132.148.132? What’s your Ethernet address?

2. Reply: I do. My Ethernet address is 00-60-8C-41-37-52

3. Ethernet Frame

   Ethernet address = 00-60-8C-41-37-52
ARP – Address Resolution Protocol

Case 2

1. Broadcast: Who has got IP address 158.132.148.132? What’s your Ethernet address?

2. Reply: The IP you indicated is not in your network. You can give the packet to me first. My MAC address is 00-60-8C-12-34-56

3. Ethernet Frame: Ethernet address = 00-60-8C-12-34-56
ARP Cache

• Will have a **heavy traffic** if so many ARP broadcast messages are generated
• Each host will have a **cache** to store the mappings (from IP to MAC address) that were obtained before

<table>
<thead>
<tr>
<th>IP Address</th>
<th>MAC Address</th>
</tr>
</thead>
<tbody>
<tr>
<td>158.132.148.80</td>
<td>00-60-8C-27-35-9A</td>
</tr>
<tr>
<td>158.132.148.28</td>
<td>02-60-8C-1A-37-49</td>
</tr>
</tbody>
</table>

• An entry will only be kept in the cache for a limited amount of time (say, 2 minutes)